Counterexample Guided Program Repair Using Large Language Models and MaxSAT-based Fault Localization

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$$P_f = ((P_o \setminus S_1) \cup S_2)$$

and

$$\forall (t_{in}^i, t_{out}^i) \in T : P_f(t_{in}^i) = t_{out}^i$$

```
int main(){ //finds max of 3 nums
       int f,s,t;
2
       scanf("%d%d%d",&f,&s,&t);
3
       if (f < s && f >= t)
          printf("%d",f);
5
       else if (s > f \&\& s >= t)
6
         printf("%d",s);
       else if (t < f \&\& t < s)
8
          printf("%d",t);
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       return 0:
11
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Symbolic-based APR Tools:

 Automated Reasoning-based tools cannot fix this program witin 90s;

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- CLARA takes too long to compute a 'minimal' repair;
- Verifix returns a compilation error.

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LLMs for code (LLMCs)

 GRANITE and CODEGEMMA cannot fix the buggy program within 90 secs;

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- GRANITE and CODEGEMMA cannot fix the buggy program within 90 secs;
- Even if we provide this assignment's description and IO tests.

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LLMs for code (LLMCs)

- GRANITE and CODEGEMMA cannot fix the buggy program within 90 secs;
- Even if we provide this assignment's description and IO tests.
- Suggesting a correct implementation, both LLMs copy the correct code, ignoring instructions not to do so.

```
int main(){ //finds max of 3 nums
       int f,s,t;
2
       scanf("%d%d%d",&f,&s,&t);
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       if (f < s && f >= t)
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 Symbolic approaches demand an excessive amount of time to produce an answer;

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```

- Symbolic approaches demand an excessive amount of time to produce an answer;
- LLMs, while fast, often produce incorrect fixes.

Program Sketches

```
int main(){ //finds max of 3 nums
       int f,s,t;
2
       scanf("%d%d%d", &f, &s, &t);
3
       if (f < g \&\& f >= t)
4
          printf("%d",f);
5
       else if (s > f \&\& s >= t)
6
          printf("%d",s);
       else if (t < f \&\& t < s)
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```

1: Program sketch with holes.

```
int main(){
       int f,s,t;
       scanf("%d%d%d",&f,&s,&t);
      @ HOLE 1 @
          printf("%d",f):
      else if (s > f \&\& s >= t)
          printf("%d",s);
      @ HOLE 2 @
          printf("%d",t);
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      return 0;
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```

Our Work

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- Combines the strengths of Formal Methods (FM) and LLM-based approaches;
- Uses MaxSAT-based fault localization to rigorously identify buggy lines, which can then be highlighted in the LLM prompt;
- For instance, instructing both LLMs to complete the previous sketch allows them to fix the buggy program in a single interaction;

MaxSAT-Based Fault Localization

• **FM24** - CFAULTS: Model-Based Diagnosis for Fault Localization in C with Multiple Test Cases.

Fault Localization - Motivation

 Debugging is one of the most time-consuming and expensive tasks in software development.

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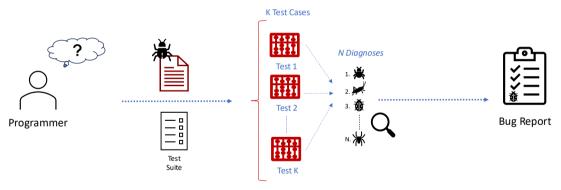
- Debugging is one of the most time-consuming and expensive tasks in software development.
- In 2024, the estimated global cost of Crowdstrike's error that hit Microsoft systems, is 5.4 Billion US\$ [The Guardian UK, 2024].

Fault Localization

• Given a buggy program, fault localization (FL) involves identifying locations in the program that could cause a faulty behaviour (bug).



 FBFL methods encode the localization problem into several optimization problems to identify a minimal set of bugs (diagnoses).



Formula-Based Fault Localization

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- do not ensure a minimal diagnosis across all failing tests (e.g., BugAssist);
- may produce an overwhelming number of **redundant diagnoses** (e.g., SNIPER).

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- We leverage MaxSAT and the theory of Model-Based Diagnosis
 (MBD) [Reiter et al., 1987], integrating all failing test cases simultaneously;
- We implement this MBD approach in a publicly available tool called CFAULTS.

Partial Maximum Satisfiability (MaxSAT)

Hard: $h_1: (v_1 \vee v_2)$ $h_2: (\neg v_2 \vee v_3)$

Soft: $s_1 : (\neg v_1)$ $s_2 : (\neg v_3)$

Partial Maximum Satisfiability (MaxSAT)

```
Hard: h_1:(v_1\vee v_2) h_2:(\neg v_2\vee v_3) Soft: s_1:(\neg v_1) s_2:(\neg v_3) Possible Solution: \neg v_1\vee v_2\vee v_3 \cos t=1
```

Figure 1: Example of a partial MaxSAT formula.

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- Each component in C can be declared **healthy** or **unhealthy**.
- For each component $c \in \mathcal{C}$, h(c) = 0 if c is unhealthy, otherwise, h(c) = 1.
- \mathcal{P} is described by a CNF formula, where \mathcal{F}_c denotes the encoding of component c:

$$\mathcal{P} \triangleq \bigwedge_{c \in \mathcal{C}} (h(c) \implies \mathcal{F}_c)$$

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- In our work, the failing test cases represent the set of observations.
- A system \mathcal{P} is considered faulty if there exists an inconsistency with a given observation o when all components are declared healthy:

$$\mathcal{P} \wedge o \wedge \bigwedge_{c \in \mathcal{C}} h(c) \vDash \bot$$

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 - is **redundant** if it is not subset-minimal [Ignatiev et al., 2019].

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To encode the MBD problem with one observation with partial MaxSAT:

- The set of clauses that encode \mathcal{P} represents the set of hard clauses;
- The soft clauses consists of unit clauses that aim to maximize the set of healthy components, i.e.,:

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- The soft clauses consists of unit clauses that aim to maximize the set of healthy components, i.e.,:

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 This encoding enables enumerating subset minimal diagnoses, considering a single observation;

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- The hard clauses, ϕ_h , in our MaxSAT formulation correspond to:

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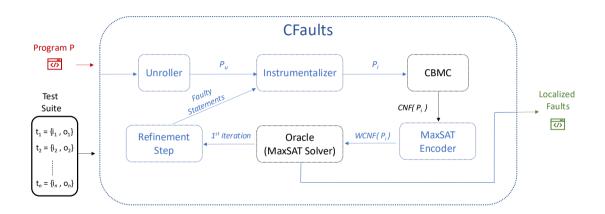
The soft clauses are formulated as:

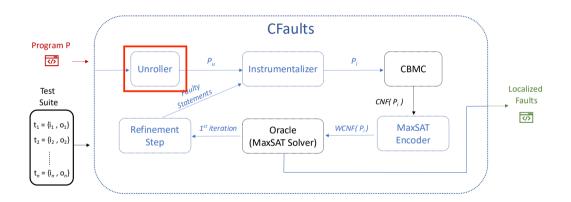
$$\phi_s = \bigwedge_{c \in \mathcal{C}} h(c).$$

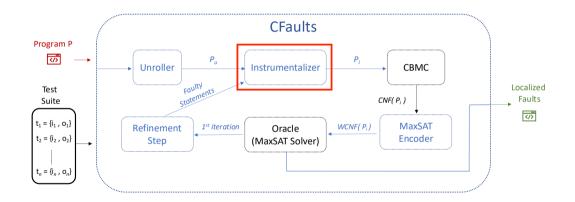
• Given a MaxSAT solution, the set of unhealthy components (h(c) = 0), corresponds to a subset-minimal aggregated diagnosis.

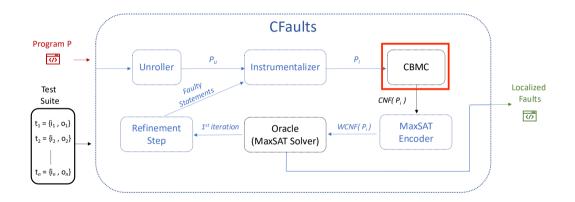
- Given a MaxSAT solution, the set of unhealthy components (h(c) = 0), corresponds to a subset-minimal aggregated diagnosis.
- This diagnosis makes the system consistent with all observations, as follows:

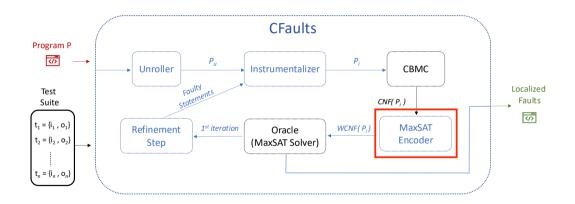
$$\bigwedge\nolimits_{o_i \in \mathcal{O}} \left(\mathcal{P}_i \wedge o_i \right) \wedge \bigwedge\nolimits_{c \in \mathcal{C} \setminus \Delta} h(c) \wedge \bigwedge\nolimits_{c \in \Delta} \neg h(c) \nvDash \bot$$

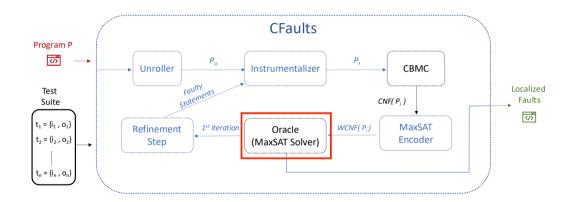








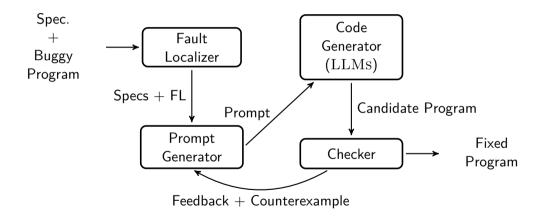




Counterexample Guided Program Repair

• AAAI 2025 - Counterexample Guided APR Using MaxSAT-based Fault Localization.

Counterexample Guided Automated Repair



Prompt Example without Fault Localization

```
# Reference Implementation
Fix all semantic bugs in the buggy program
                                              (Do not copy this program) <c> #
below. Modify the code as little as possible.
                                              ...c
Do not provide any explanation.
                                              int main(){
                                                // Reference Implementation
### Problem Description ###
Write a program that determines and
                                              . . .
prints the largest of three integers
given by the user.
                                              ### Buggy Program <c> ###
                                              ...
### Test Suite
                                              int main(){
#input:
                                                // Buggy program from Listing 1
6 2 1
#output:
                                              . . .
6
// The other input-output tests
                                              ### Fixed Program <c> ###
                                              ```c
```

# **Prompt with Fault Localization (Sketches)**

```
Complete all the '@ HOLES N @' in the
incomplete program below.
Modify the code as little as possible.
Do not provide any explanation.
Problem Description
Write a program that determines and
prints the largest of three integers
given by the user.
Test Suite
#input:
6 2 1
#output:
6
// The other input-output tests
```

```
Reference Implementation
(Do not copy this program) <c> #
```c
int main(){
 // Reference Implementation
### Incomplete Program <c> ###
```c
int main(){
 // Buggy program from Listing 1
Complete Program <c>
```c
```

Experimental Results

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 - Microsoft's Рні3.
- Experiments were conducted using a memory limit of 10GB, and a timeout of 90s.

LLM-Driven APR with CFaults

LLMs	Prompt Configurations				
	De-TS	De-TS-CE	Sk_De-TS	Sk_De-TS-CE	Portfolio (All Configurations)
CodeGemma	597 (41.7%)	606 (42.3%)	682 (47.7%)	688 (48.1%)	823 (57.5%)
CodeLlama	492 (34.4%)	500 (34.9%)	573 (40.0%)	561 (39.2%)	712 (49.8%)
Gemma	496 (34.7%)	492 (34.4%)	532 (37.2%)	534 (37.3%)	670 (46.8%)
Granite	626 (43.7%)	624 (43.6%)	691 (48.3%)	681 (47.6%)	846 (59.1%)
Llama3	564 (39.4%)	590 (41.2%)	578 (40.4%)	591 (41.3%)	851 (59.5%)
Phi3	494 (34.5%)	489 (34.2%)	547 (38.2%)	535 (37.4%)	621 (43.4%)
Portfolio (All LLMs)	842 (58.8%)	846 (59.1%)	900 (62.9%)	907 (63.4%)	1013 (70.8%)
Verifix	90 (6.3%)				
Clara	495 (34.6%)				

Table 1: The number of programs fixed by each LLM under various configurations. Mapping abbreviations to configuration names: **De** - IPA *Description*, **TS** - *Test Suite*, **CE** - *Counterexample*, **SK** - *Sketches*.

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- All six LLMs using different prompt configurations repair more programs than traditional APR tools;
- Incorporating FL-based Sketches allows LLMs to repair more programs;
- Including a reference implementation allows for more repaired programs but with less efficient fixes (see our paper);
- Our CEGIS approach significantly improves the accuracy of LLM-driven APR across various configurations;

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- We employ MaxSAT-based Fault Localization to guide and minimize LLMs' patches to incorrect programs by feeding them bug-free program sketches;

- We tackle the APR problem using an LLM-Driven Counterexample Guided Inductive Synthesis (CEGIS) approach [Solar-Lezama et al., 2006];
- We employ MaxSAT-based Fault Localization to guide and minimize LLMs' patches to incorrect programs by feeding them bug-free program sketches;
- With our approach all six evaluated LLMs fix more programs and produce smaller patches than other configurations and symbolic tools;

- We tackle the APR problem using an LLM-Driven Counterexample Guided Inductive Synthesis (CEGIS) approach [Solar-Lezama et al., 2006];
- We employ MaxSAT-based Fault Localization to guide and minimize LLMs' patches to incorrect programs by feeding them bug-free program sketches;
- With our approach all six evaluated LLMs fix more programs and produce smaller patches than other configurations and symbolic tools;
- Our code is available on GitHub and on Zenodo.

Thank you!



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